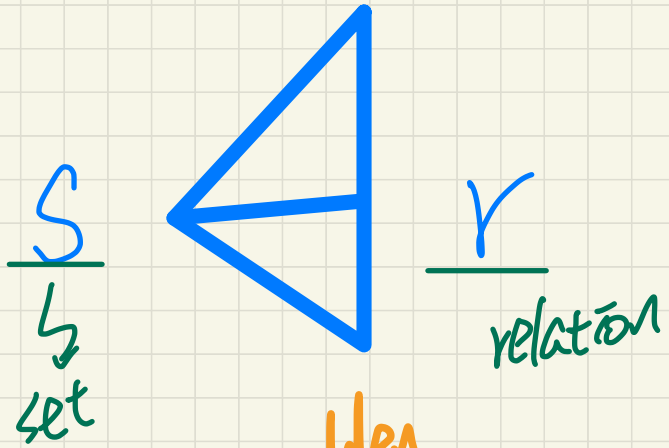
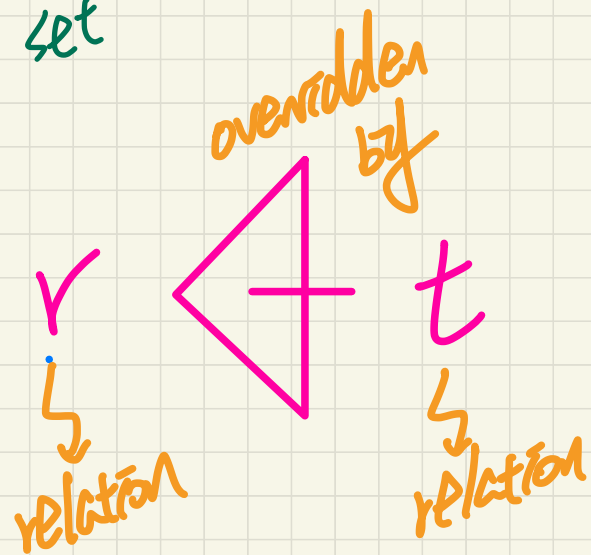


Lecture 1b

Review on Math (continued)



domain subtraction



overriding



$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$



Definition: $r \triangleleft t = \{ (d, r) \mid (d, r) \in t \vee ((d, r) \in r \wedge d \notin \text{dom}(t)) \}$
 e.g., $r \triangleleft \{(a, 3), (c, 4)\}$

union

relation

another relation

$$r \triangleleft \{(a, 3), (c, 4)\}$$

$$\text{dom}(t) = \{a, c\}$$

$$\{(a, 3), (c, 4)\} \cup \{(b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\}$$

$$= \{ \underline{\hspace{10em}} \}$$

$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

$$r[s] = \text{ran}(\underset{\textcircled{s}}{s} \triangleleft r)$$

$$r[\underbrace{\{a, b\}}_{\textcircled{s}}] = \text{ran}(\underbrace{\{a, b\}}_{\textcircled{s}} \triangleleft r)$$

$$= \text{ran}(\{(a, 1), (b, 2), (a, 4), (b, 5)\})$$

$$= \{1, 2, 4, 5\}$$

Side Note -
databases
↳ relational databases
(SQL queries)
↳ relational algebra

$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$

$$r \triangleleft t = t \cup (\text{dom}(t) \triangleleft r)$$

algebraic property.

$$a + b = b + a$$

$$r \triangleleft \underbrace{\{(a, 3), (c, 4)\}}_t$$

$$r \triangleleft \underbrace{\{(a, 3), (c, 4)\}}_t$$

$$= \{(a, 3), (c, 4)\} \cup (\{a, c\} \triangleleft r)$$

$$\{(b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\}$$

$$= \{ \underline{\hspace{10em}} \}$$

$isFunctional(r)$

\iff

$\forall s, t_1, t_2 \bullet \underline{(s \in S \wedge t_1 \in T \wedge t_2 \in T)} \Rightarrow ((s, t_1) \in r \wedge (s, t_2) \in r \Rightarrow t_1 = t_2)$

to disprove \rightarrow
find witness
(satisfying antecedent
but violating consequent)

||| Contrapositive $P \Rightarrow Q \equiv \neg Q \Rightarrow \neg P$

$t_1 \neq t_2 \Rightarrow (s, t_1) \notin r \vee (s, t_2) \notin r$

What is the smallest relation satisfying the functional property?

$\hookrightarrow \emptyset$ \because we cannot find a witness to disprove that it violates the functional property $\therefore \emptyset$ is a function

$\text{dom} = \{z \rightarrow 1\} \subseteq \mathbb{C}$

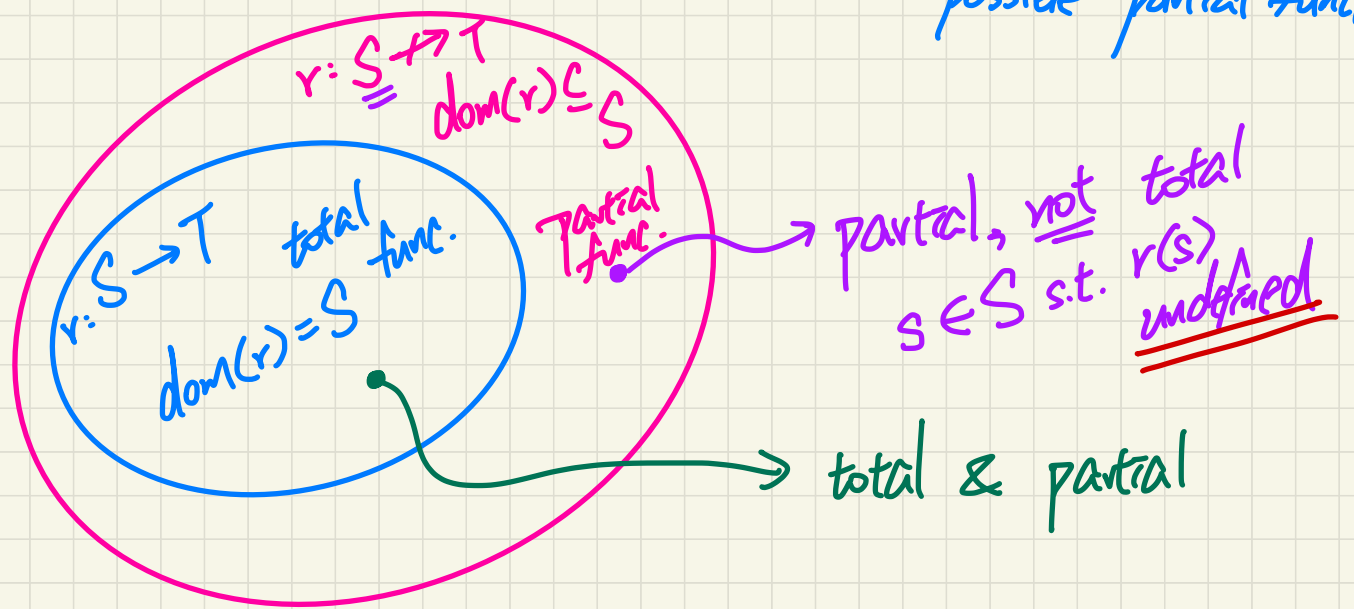
 $\text{dom} = \{z \rightarrow 3, 1\} \subseteq \mathbb{C}$

e.g., $\{ \{(2, a), (1, b)\}, \{(2, a), (3, a), (1, b)\} \} \subseteq \{1, 2, 3\} \rightarrow \{a, b\}$

function

function

the set of possible partial functions



	injective	surjective	bijective
partial	.	.	.
total	.	.	.

Injective Functions

$isInjective(f)$

\iff

$$\forall s_1, s_2, t \bullet (s_1 \in S \wedge s_2 \in S \wedge t \in T) \Rightarrow ((s_1, t) \in f \wedge (s_2, t) \in f \Rightarrow s_1 = s_2)$$

$b = b \Rightarrow I = 3$ (False)

If f is a **partial injection**, we write: $f \in S \rightsquigarrow T$

- e.g., $\{\emptyset, \{(1, a)\}, \{(2, a), (3, b)\}\} \subseteq \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, b), (2, a), (3, b)\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, b), (3, b)\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$

$\not\rightarrow \emptyset$ \neq ! not total
 2. injective
 \therefore no witnesses of violation

If f is a **total injection**, we write: $f \in S \rightarrow T$

- e.g., $\{1, 2, 3\} \rightarrow \{a, b\} = \emptyset \rightarrow \{(1, a), (2, b), (3, a)\}$
- e.g., $\{(2, d), (1, a), (3, c)\} \in \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c), (3, d)\} \in \{1, 2, 3\} \rightarrow \{a, b, c, d\}$

\rightarrow
 not total, inj. (False)
 total, not inj. $(2, d), (3, d)$ $d = d \Rightarrow 2 = 3$

partial, not inj.

the set of all possible total injections

Surjective Functions

$$isSurjective(f) \iff \underline{ran}(f) = \underline{T}$$

If f is a **partial surjection**, we write: $f \in S \dashrightarrow T$

- e.g., $\{(1, b), (2, a)\}, \{(1, b), (2, a), (3, b)\} \subseteq \{1, 2, 3\} \dashrightarrow \{a, b\}$
- e.g., $\{(2, a), (1, a), (3, a)\} \not\subseteq \{1, 2, 3\} \dashrightarrow \{a, b\}$ $ran = \{a\}$ *partial, not sur.*
- e.g., $\{(2, b), (1, b)\} \not\subseteq \{1, 2, 3\} \dashrightarrow \{a, b\}$ $ran = \{b\}$ *partial, not sur.*

\dashrightarrow

If f is a **total surjection**, we write: $f \in S \rightarrow T$

- e.g., $\{(2, a), (1, b), (3, a)\}, \{(2, b), (1, a), (3, b)\} \subseteq \{1, 2, 3\} \rightarrow \{a, b\}$
- e.g., $\{(2, a), (3, b)\} \not\subseteq \{1, 2, 3\} \rightarrow \{a, b\}$ $dom = \{2, 3\}$ *not total, sur.*
- e.g., $\{(2, a), (3, a), (1, a)\} \not\subseteq \{1, 2, 3\} \rightarrow \{a, b\}$ $ran = \{a\}$

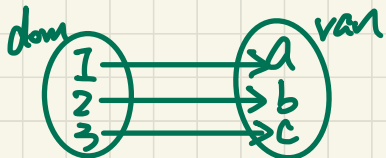
\rightarrow

total, sur.
sur.

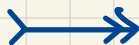
not sur.

total, not sur.

Bijjective Functions



f is **bijjective**/a **bijection**/one-to-one correspondence if f is **total**, **injective**, and **surjective**.



- e.g., $\{1, 2, 3\} \mapsto \{a, b\} = \emptyset \quad \{(1, a), (2, b), (3, ?)\}$
- e.g., $\{(1, a), (2, b), (3, c)\}, \{(2, a), (3, b), (1, c)\} \subseteq \{1, 2, 3\} \mapsto \{a, b, c\}$
- e.g., $\{(2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \mapsto \{a, b, c\}$ **not total, inj, sur.**
- e.g., $\{(1, a), (2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \mapsto \{a, b, c\}$ **total, not inj, sur.**
- e.g., $\{(1, a), (2, c)\} \notin \{1, 2\} \mapsto \{a, b, c\}$ **ran = $\{a, c\}$**

total ✓

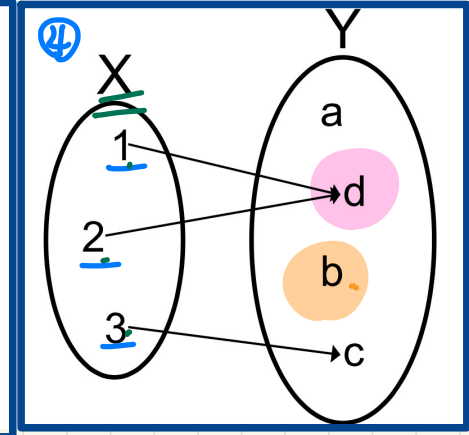
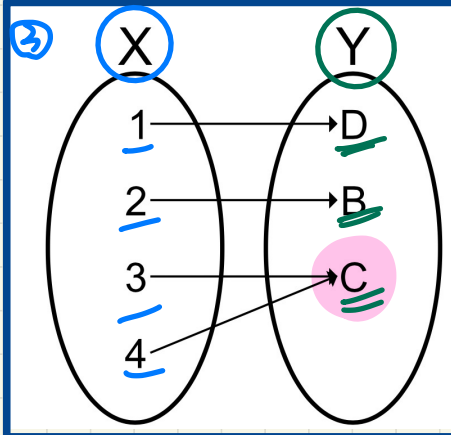
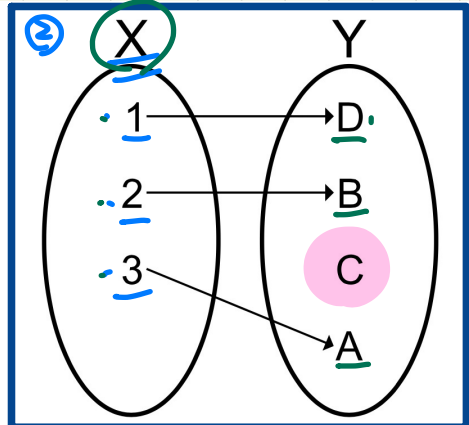
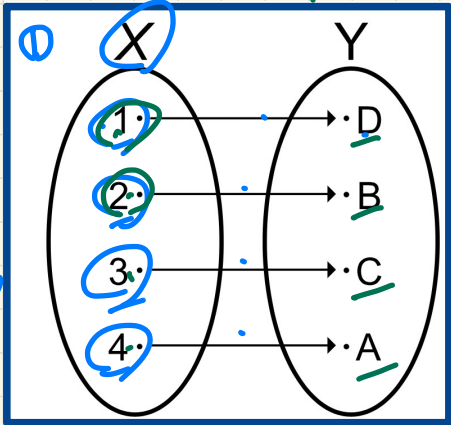
inj. ✓

sur. ✗

Exercise

$\text{dom}(\mathbb{1}) = X$
 $\text{ran}(\mathbb{1}) = Y$

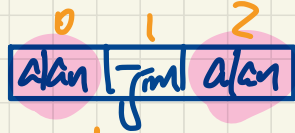
✓
Exercise
 Make a
 function that's
 partial but
 not total.



	①	②	③	④
partial	✓	✓	✓	✓
total	✓	✓	✓	✓
inj.	✓	✓	✗	✗
sur.	✓	✗	✓	✗
bij.	✓	✗	✗	✗

Formalizing Arrays as Functions

Not partial inj.

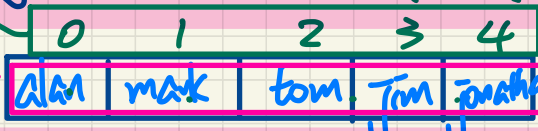


indices \approx domain

```
String[] a = new String[5];
```

$a = \{(0, alan), (1, jim), (2, alan)\}$
programming

content range \approx



Strings 0 "" 2 "" 3 5

$$a = \{(0, "alan"), (1, "mark"), (2, "tom"), (3, "jim"), (4, "jordan")\}$$

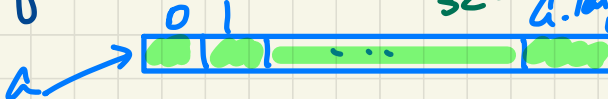
formalization in math.

Should a be formalized/modelled as a relation?

No. $\because \{(0, alan), (0, jim)\}$

$\mathbb{Z} \leftrightarrow \text{String}$

Partial Injection



in reality, only one element may be stored at each index